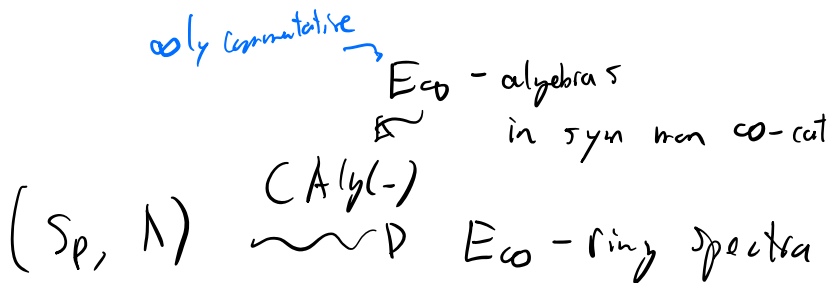
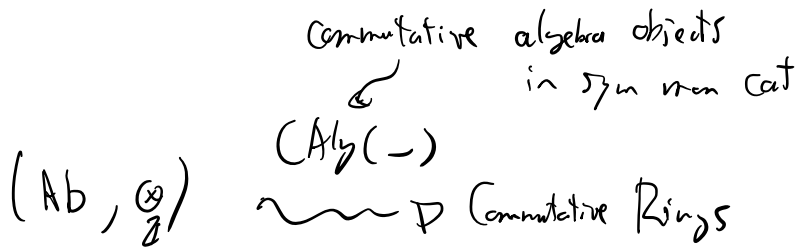


Sym mon co-cats



Sym mon 1-cats recall

A lax sym mon functor $F: \mathcal{C} \rightarrow \mathcal{D}$ has

natural transformations $F(c) \otimes F(c') \rightarrow F(c \otimes c')$, $1 \rightarrow F(1e)$

Compatibility with α & β
assoc prod

Ex an algebra in \mathcal{C} is a lax mon functor $* \rightarrow \mathcal{C}$

terminal (unique sym mon str on a pt)

Need to encode ∞ many compatibilities such as pentagon 'cocycle homotopy'

Def $Fin_{\neq} := \{ \text{finite pointed sets} \} \ni \{ \neq, 1, \dots, n \} =: \langle n \rangle$

Combinatorics machine Γ^{op}

$m_n: \langle n \rangle \rightarrow \langle 1 \rangle$ $m_n(i) = 1$ if $i \neq *$

"operadic stuff"

$g_i: \langle n \rangle \rightarrow \langle 1 \rangle$ insert \leftarrow throw shit away

$j \mapsto \begin{cases} 1 & j=i \\ * & j \neq i \end{cases}$

$|a^{-1}(i)| = 1 \quad \forall i \neq *$

Segal's special Γ -spaces: $(\text{co}1)\text{-cat}$

Def'n A sym mon co-cat is a co-functor

$\underline{C} : \text{Fin}_* \rightarrow \text{Cat}_\infty$ satisfying the Segal condition:

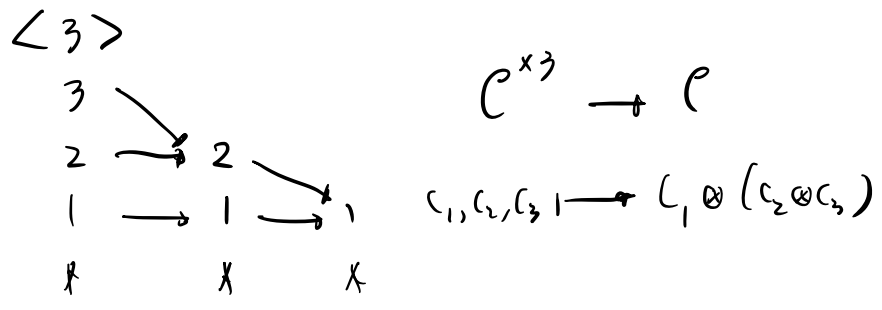
s.t. $\forall n \geq 0$ $\underline{C}(\langle n \rangle) \xrightarrow{(\text{g}^i)} \prod_{i=1}^n \underline{C}(\langle 1 \rangle)$ equivalence

e.g. $n=0 \Rightarrow \underline{C}(\langle 0 \rangle) = *$

choice of inverse needed for def'n of \otimes

$\underline{C} := \underline{C}(\langle 1 \rangle) \quad \otimes : \underline{C} \times \underline{C} \xrightarrow{(\text{g}^i, \text{g}^j)} \underline{C}(\langle 2 \rangle) \xrightarrow{m_2} \underline{C}(\langle 1 \rangle)$

$\mathbb{1} : * = \underline{C}(\langle 0 \rangle) \rightarrow \underline{C}(\langle 1 \rangle) = \underline{C}$



writer, braiding

with counit can derive pentagon (which is a nat. modification)

Claim sym mon 1-cat $\xRightarrow{\text{Nerve}}$ sym mon co-cat

$\text{Fin}_* \rightarrow \text{Cat}_\infty$
 $\langle n \rangle \mapsto N(\mathcal{C})^{\times n}$

def'n A strong symmetric functor is a nat transform $\underline{C} \Rightarrow \underline{D}$

$\Rightarrow \text{Sym Mon Cat}_0 \subseteq \text{Fun}(\text{Fin}_n, \text{Cat}_0)$ full subset

problem We are considering Cat_0 as a $(\infty, 1)$ cat, not $(\infty, 2)$ cat.

$(\infty, 2)$ is necessary for lex sym mon functors

Prop \mathcal{P} sym mon \Rightarrow h \mathcal{P} sym mon (converse false)

Remark Fin_n and Δ^{op} for monoidal

\uparrow
E ω

\nwarrow
A ω

determines combinatorics

Grothendieck op-fibrations

first 1-categorical: $F: \mathcal{B} \rightarrow \text{Cat}_1$ ^{kind of c-cat} 2-functor pseudo functor

Def the Grothendieck construction of $F: \mathcal{B} \rightarrow \text{Cat}_1$ category of elements

has obj: (b, e) $b \in \text{ob } \mathcal{B}, e \in \text{ob } F(b)$

mor $(b_1, e) \rightarrow (b_2, e): \alpha: b_1 \rightarrow b_2 \quad \exists: F(\alpha)(e) \rightarrow e_2$

\exists forgetful functor $\int F \xrightarrow{p} \mathcal{B} \quad (b, x) \mapsto b \quad p^{-1}(b) = F(b)$
associated grothendieck fibration

Th $\text{Fun}(\mathcal{B}, \text{Cat}_1) \rightarrow \text{Cat}_1 / \mathcal{B}$
 $F \mapsto \left(\int F \rightarrow \mathcal{B} \right)$ ^{fibration}

is faithful onto the (non-full) subcategory of $\text{Cat}_1 / \mathcal{B}$ with

Grothendieck / criterion of fibration as objects, "criterion functors" as 1-morphisms
(fibered categories)

& "vertical natural transformations" as 2-morphisms

\tilde{c} takes cartesian arrows to cartesian arrows

coExample stacks $\text{Mod } M \rightarrow \text{Cat}_1, M \mapsto \text{Vect } M$ (not op)

Example G group, $F: \mathcal{B} \rightarrow \text{1-cat}$ ^{discrete} any 2-functor $\mathcal{P} := F(x)$
 \mathcal{B} 1-cat with 1 object & G as morphisms
"action $G \curvearrowright \mathcal{P}$ "

$\text{ob } \int F = \text{ob } \mathcal{P}$

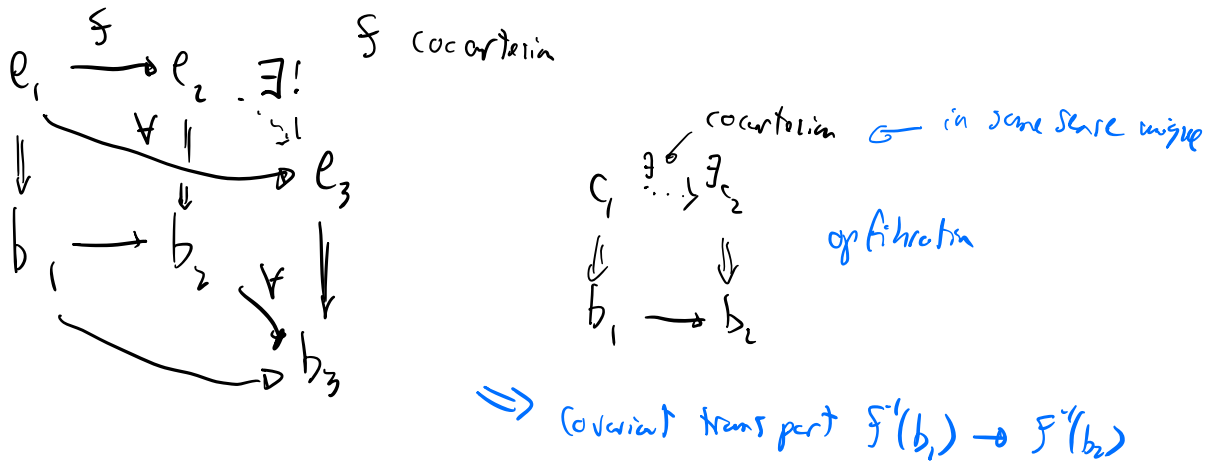
$\text{mor}_{\int F}(d_1, d_2) = \{ \alpha = g \in G, F(g)[d_1] \rightarrow d_2 \}$

composition uses $F(g_1) \circ F(g_2) \simeq F(g_1 g_2)$ "homology quotient"

subexamp $\mathcal{C} = *//H \Leftrightarrow F$ "nonabelian cocycle"

$$\& \quad 1 \rightarrow H \rightarrow \text{mor} \downarrow F \rightarrow G \rightarrow 1 \quad \text{(classifies extensions)}$$

\uparrow
fiber
 \uparrow
forgetful functor



(\mathcal{C}, \otimes) sym mon cat $\Rightarrow F: \text{Fin}_* \rightarrow \text{Cat}_\infty$

the category of operators is $\int F$:

obj: possibly empty finite ordered lists $c_1, \dots, c_n \in \mathcal{C}$

mor $((c_1, \dots, c_n), (d_1, \dots, d_m))$

consists of $\alpha: \langle n \rangle \rightarrow \langle m \rangle \in \text{Fin}_*$ & $\forall k \in \langle m \rangle \setminus \text{Ext}$

a map \otimes

$$\nearrow \text{id}^{-1}(\langle k \rangle) \quad c_i \rightarrow d_k$$

$$c_1 \otimes c_2 \otimes c_3 \Rightarrow (c_1 \otimes c_2) \otimes c_3$$

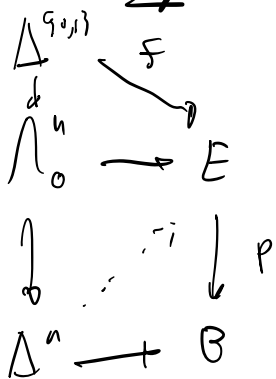
obvious composition

forgetful

$$\mathcal{C}^\otimes \xrightarrow{p} \text{Fin}_*$$

Def $p: E \rightarrow B$ \in mod str

An edge of E is p -cocartesian if every lifting problem



has a solution $\forall n \geq 2$

several equivalent def'ns

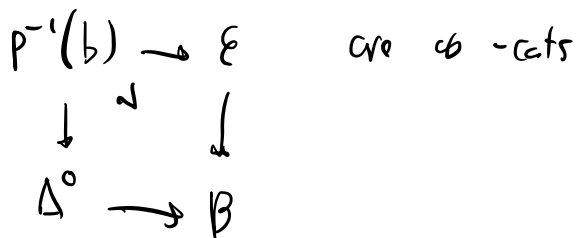
p is cocartesian fibration if

(i) p inner fibration (Fibers in Joyal model str) \leftarrow fibrant = co-cat

(ii) \forall edge $\alpha: b_1 \rightarrow b_2$ in B $\forall e_1 \in p^{-1}(b_1)$ vertex \leftarrow

\exists p -cocartesian edge $f: e_1 \rightarrow e_2$ s.t. $p(f) = \alpha$

The fibers



Prop For E, B co-cats, fibration

assemble into a functor $\alpha: b_1 \rightarrow b_2 \rightsquigarrow \alpha_!: p^{-1}(b_1) \rightarrow p^{-1}(b_2)$

"straightening" Grothendieck construction = "unstraightening"

Def'n of Lurie A symmetric co-cat is a cocartesian fibration

$$C^{\otimes} \rightarrow \text{Fin}_*$$

$$\left(\begin{array}{c}
 \mathcal{C}_{\langle n \rangle}^{\otimes} \rightarrow \mathcal{C}^{\otimes} \\
 \downarrow \quad \downarrow \\
 \Delta^0 \xrightarrow{\langle n \rangle} \text{Fin}_n \\
 \& \exists \text{ lift } \alpha_! : \mathcal{C}_{\langle n \rangle}^{\otimes} \rightarrow \mathcal{C}_{\langle m \rangle}^{\otimes} \quad \forall \alpha : \langle n \rangle \rightarrow \langle m \rangle \in \text{Fin}_* \\
 \& \text{ the maps } \mathcal{C}_{\langle n \rangle}^{\otimes} \xrightarrow{(\beta_i)_!} \mathcal{C} := \mathcal{C}_{\langle 1 \rangle}^{\otimes} \\
 \text{are equivalences } \forall n \geq 0
 \end{array} \right)$$

Example \mathcal{C} with finite products & $\otimes = \times$

for experts
 Remark \downarrow
 N_{Δ} (nice simplicial model cat with $\text{sym} \otimes$) is sym mon
 \uparrow
 simplicial nerve

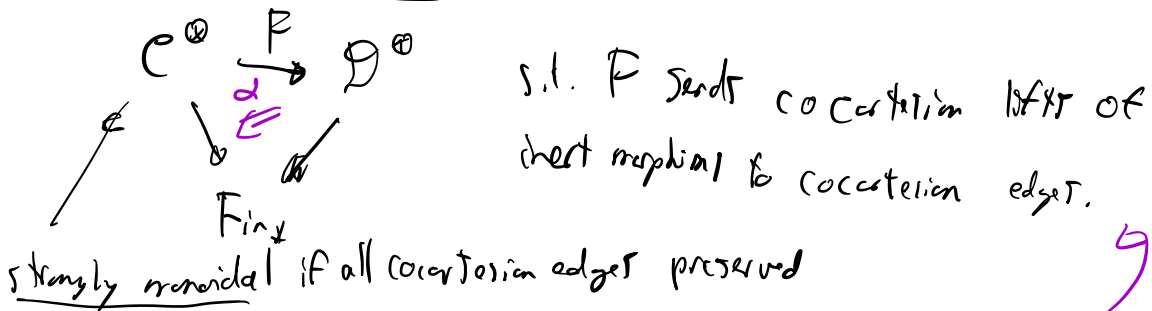
The Sp & $\mathcal{D}(\mathbb{R})$ admit unique strict closed
 sym mon ω -cat with units \mathbb{J} & $\mathbb{R}[0]$

$$\otimes = \wedge \quad \& \quad \otimes = \otimes_{\mathbb{R}}^{\mathbb{L}}$$

$-\otimes_{\mathcal{C}} : \mathcal{C} \rightarrow \mathcal{C}$ admits a right adjoint $\text{Hom}_{\mathcal{C}}(_, -)$

$(\text{Ch}^{-1} \text{cat})$ ["quasi-invert"] $\simeq \mathcal{D}^{-1} \text{cat}$ or dg-Nerve of fib-cat objects

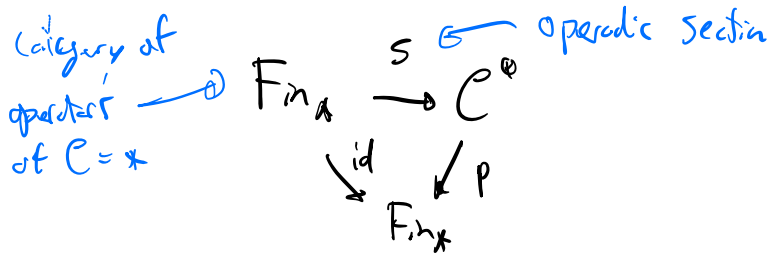
Def A lax symmetric monoidal functor is a functor F and α



Note α is unique if it exists, because of the conditions

2. A lax symmetric monoidal functor of symmetric \mathcal{C} -cat induces one of \mathcal{C} -cats

Def An algebra in \mathcal{C}^{\otimes} is a lax symmetric $\ast \rightarrow \mathcal{C}^{\otimes}$



Concretely: $A := s\langle 1 \rangle \in \mathcal{C}$ "the actual algebra"

$$s\langle 2 \rangle = (A_1, A_2) \in \mathcal{C}\langle 2 \rangle \cong \mathcal{C} \times \mathcal{C}$$

$$\& s\langle 2 \rangle \xrightarrow{s(m_2)} s\langle 1 \rangle \text{ gives } A_1 \otimes A_2 \rightarrow A$$

$$s(\beta_i) \text{ cocartesian} \Rightarrow \text{supplies } A_1 \cong A_2 \cong A$$

$$s\langle 0 \rangle \in \mathcal{C}\langle 0 \rangle \rightarrow \mathcal{C}\langle 1 \rangle \text{ gives unit}$$

Commutativity, ...

Def A E_0 -ring spectrum is an algebra object in (Sp, \wedge)

Can now define moduli, limits, spectrum objects in \mathcal{C}

$\mathcal{C} \ni \mathcal{S} \xrightarrow{(\cdot)_*} \mathcal{S}_* \xrightarrow{\Sigma^0} Sp$ \mathcal{S} in non functors